

# Communications protocol implications of using code-based cryptography

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NCCoE workshop on migrating to post-quantum cryptography

# History of code-based cryptography

1978 McEliece: Public-key encryption using error-correcting codes.

- ▶ Original parameters designed for  $2^{64}$  security.  
Using 1962 Prange information-set decoding for parameter choice.
- ▶ 2008 Bernstein–Lange–Peters: broken original parameters in  $\approx 2^{60}$  cycles.
- ▶ Easily scale up for higher security.
- ▶ The McEliece system (with later key-size optimizations) achieves  $2^\lambda$  security against Prange's attack using  $(0.741186 \dots + o(1))\lambda^2(\log_2 \lambda)^2$ -bit keys as  $\lambda \rightarrow \infty$ .

## Security analysis of McEliece encryption

Some papers studying algorithms for attackers:

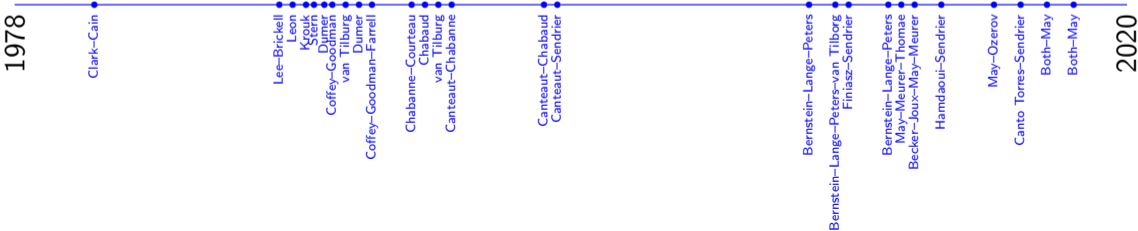
1962 Prange; 1981 Clark–Cain, crediting Omura; 1988 Lee–Brickell; 1988 Leon; 1989 Krouk; 1989 Stern; 1989 Dumer; 1990 Coffey–Goodman; 1990 van Tilburg; 1991 Dumer; 1991 Coffey–Goodman–Farrell; 1993 Chabanne–Courteau; 1993 Chabaud; 1994 van Tilburg; 1994 Canteaut–Chabanne; 1998 Canteaut–Chabaud; 1998 Canteaut–Sendrier; 2008 Bernstein–Lange–Peters; 2009 Bernstein–Lange–Peters–van Tilborg; 2009 Bernstein (**post-quantum**); 2009 Finiasz–Sendrier; 2010 Bernstein–Lange–Peters; 2011 May–Meurer–Thomae; 2012 Becker–Joux–May–Meurer; 2013 Hamdaoui–Sendrier; 2015 May–Ozerov; 2016 Canto Torres–Sendrier; 2017 Kachigar–Tillich (**post-quantum**); 2017 Both–May; 2018 Both–May; 2018 Kirshanova (**post-quantum**).

All of these attacks involve huge searches, like attacking AES.

The quantum attacks (Grover etc.) leave at least half of the bits of security.

# Attack progress over time

$$\lim_{K \rightarrow \infty} \frac{\log_2 \text{AttackCost}_{\text{year}}(K)}{\log_2 \text{AttackCost}_{2020}(K)}$$



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∞

Red: Lattices have lost much more security.  
 Lattices had 42% higher security levels  
 ten years ago than they have today.

1978

Clark-Gain

Lee-Brickell  
 Leon  
 Krouk  
 Stern  
 Dumer  
 Coffey-Goodman  
 van Tilburg  
 Dumer  
 Coffey-Goodman-Farrell  
 Chabanne-Courteau  
 Chabaud  
 van Tilburg  
 Canteaut-Chabanne

Canteaut-Chabaud  
 Canteaut-Sendrier

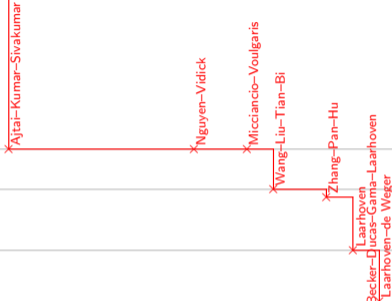
Bernstein-Lange-Peters  
 Bernstein-Lange-Peters-van Tilburg  
 Fimiasz-Sendrier

Bernstein-Lange-Peters  
 May-Meurer-Thomae  
 Becker-Joux-May-Meurer  
 Hamdaoui-Sendrier

May-Ozerov  
 Becker-Ducas-Gama-Laarhoven  
 Laarhoven-de Weger  
 Canto Torres-Sendrier  
 Both-May  
 Both-May

2020

1.421  
 1.315  
 1.154



# NIST PQC submission Classic McEliece

No patents. ✓

Shortest ciphertexts of all Round-2 candidates. ✓

Fast open-source constant-time software implementations. ✓

Very conservative system, expected to last; has strongest security track record. ✓

Sizes with similar post-quantum security to AES-128, AES-192, AES-256:

<b>Metric</b>	<b>mceliece348864</b>	<b>mceliece460896</b>	<b>mceliece6960119</b>
Public-key size	261120 bytes	524160 bytes	1047319 bytes
Secret-key size	6452 bytes	13568 bytes	13908 bytes
Ciphertext size	128 bytes	188 bytes	226 bytes
Key-generation time	52415436 cycles	181063400 cycles	417271280 cycles
Encapsulation time	43648 cycles	77380 cycles	143908 cycles
Decapsulation time	130944 cycles	267828 cycles	295628 cycles

See <https://classic.mceliece.org> for authors, details & parameters.

# Key issues for McEliece

# BIG PUBLIC KEYS.



# Key issues for McEliece

- ▶ Sending 1MB takes time and bandwidth.
- ▶ Google–Cloudflare experiment:

*in some cases the public-key + ciphertext size was too large to be viable in the context of TLS*

and even 10KB messages dropped.

But users send big data anyway. We have lots of bandwidth.  
A key takes less space than a kitten picture.



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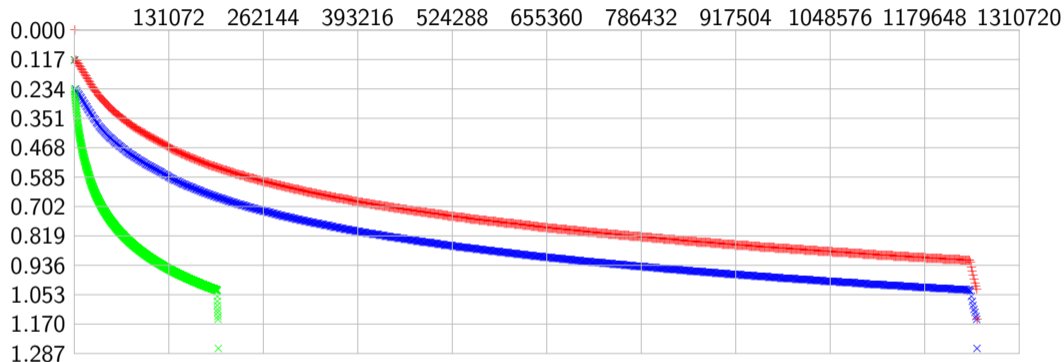
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- ▶ But: If any client is allowed to send a new ephemeral 1MB McEliece key to server, an attacker can easily flood server's memory. **This invites DoS attacks.**

# Use cases for Classic McEliece

- ▶ Standard public-key encryption, e.g. GnuPG/PGP with long-term keys.
- ▶ **PQ-Wireguard** (Hülsing, Ning, Schwabe, Weber, Zimmermann; IEEE S&P 2021).
  - ▶ Uses McEliece for long-term identity key in KEM-KEM construction.
  - ▶ McEliece key exchanged out of band at registration.
  - ▶ Strong benefit from short ciphertexts.
  - ▶ Combined with lattice-based scheme for ephemeral keys.
- ▶ **McTiny** (Bernstein, Lange; USENIX 2020)
  - ▶ McEliece also used for ephemeral keys.
  - ▶ Avoids DoS memory flooding attacks by using structure of code-based encryption. Server returns partial encryption and state in cookie encrypted to itself; cookie is smaller than network packet sent to server.
  - ▶ Good speed and security with congestion control and surrounding protocol.

## Measurements of our software (<https://mctiny.org>)



Client time vs. bytes sent, bytes acknowledged, bytes in acknowledgments.  
Curve shows packet pacing from our new user-level congestion-control library.