## Cryptology, homework sheet 2

Due: 22 September 2016, 10:45 for students of 2MMC10 and 06 October 2016, 10:45 for students following the MasterMath course.

2MMC10: Please hand in your homework in groups of two or three. To submit your homework, place it on the table of the lecturer *before* the lecture.

Mastermath: Instructions on how to submit will follow. Please team up in groups of 2 or 3.

Please write the names and student numbers on the homework sheet. Please indicate your home university and study direction.

This time one-line answers using a computer algebra system do *not* count. But it is a good moment to familiarize yourself with some system(s) so that you know how to solve similar problems for real life examples and to verify your answers. You may use a computer algebra system to compute subresults, such as f div g and  $f \cdot g$ . See below for a description of the Extended Greatest Common Divisor Algorithm (XGCD).

- 1. Compute the extended gcd of 155 and 649 using XGCD.
- 2. Compute the extended gcd of  $f(x) = x^5 + 3x^3 + x^2 + 2x + 1$  and  $g(x) = x^4 5x^3 5x^2 5x 6$  in  $\mathbb{Q}[x]$  using XGCD.
- 3. Consider the residue classes of  $\mathbb{F}_2[x]$  modulo  $f(x) = x^n + 1$  for some positive integer n > 1, i.e.  $R = \mathbb{F}_2[x]/(x^n + 1)$ . Note that R can be represented as

$$R = \left\{ a_0 + a_1 x + a_2 x^2 + \ldots + a_{n-1} x^{n-1} \mid a_i \in \mathbb{F}_2 \right\}.$$

Show that R is not a field.

Hint: Find a non-zero element that is not invertible.

4. Let K be a field of characteristic p, where p is prime. Show that for any integer  $n \ge 0$  one has

$$(a+b)^{p^n} = a^{p^n} + b^{p^n}$$

for all  $a, b \in K$ .

Hint: You can use the binomial theorem and use proof by induction.

5. Use the Rabin test (see below) to prove that  $x^4 + x + 1$  is irreducible over  $\mathbb{F}_2$ . You should be able to do this exercise by hand. Please document the results of all steps in the algorithm and show how they were obtained.

Here is a description of XGCD. This description assumes that the input elements f, g live in some ring R in which the greatest common divisor is defined. We will usually use the XGCD on integers or polynomials. If the inputs are integers you can ignore the part the leading coefficient.

## Algorithm 1 (Extended Euclidean algorithm)

IN:  $f, g \in R$ OUT:  $d, u, v \in R$  with d = uf + vg1.  $a \leftarrow [f, 1, 0]$ 2.  $b \leftarrow [g, 0, 1]$ 3. repeat (a)  $c \leftarrow a - (a[1] \operatorname{div} b[1])b$ (b)  $a \leftarrow b$ (c)  $b \leftarrow c$ while  $b[1] \neq 0$ 4.  $l \leftarrow LC(a[1]), a \leftarrow a/l /*LC = leading coefficient, this only applies to polynomials*/$  $5. <math>d \leftarrow a[1], u \leftarrow a[2], v \leftarrow a[3]$ 6. return d, u, v

In this algorithm, div denotes division with remainder. The first component of c is thus easier written as  $c[1] \leftarrow a[1] \mod b[1]$  but by operating on the whole vector we get to update the values leading to u and v, too. At each step we have

$$a[1] = a[2]f + a[3]g$$
 and  $b[1] = b[2]f + b[3]g$ .

To see this, note that this holds trivially for the initial conditions. If it holds for both a and b then also for c since it computes a linear relation of both vectors. So each update maintains the relation and eventually when b[1] = 0, we have that a[1] holds the previous remainder, which is the gcd of f and g. If the inputs are polynomials, at the end the gcd is made monic by dividing by the leading coefficient LC(a[1]).

**Example 2** Let  $R = \mathbb{R}[x]$  and  $f(x) = x^5 + 3x^3 - x^2 - 4x + 1$ ,  $g(x) = x^4 - 8x^3 + 8x^2 + 8x - 9$ . So at first we have a = [f, 1, 0], b = [g, 0, 1].

We have  $(a[1] \operatorname{div} b[1]) = x + 8$  and so end the first round with

$$a = [g, 0, 1],$$
  

$$b = [59x^3 - 73x^2 - 59x + 73, 1, -x - 8].$$

Indeed b[1] = f(x) + (-x - 8)g(x).

With these new values we have  $(a[1] \operatorname{div} b[1]) = 1/59x - 399/3481$  and so the second round ends with

$$\begin{array}{ll} a & = & [59x^3 - 73x^2 - 59x + 73, 1, -x - 8], \\ b & = & [2202/3481x^2 - 2202/3481, -1/59x + 399/3481, 1/59x^2 + 73/3481x + 289/3481]. \end{array}$$

In the third round we have  $(a[1] \operatorname{div} b[1]) = 205379/2202x - 254113/2202$  and obtain

$$a = [2202/3481x^2 - 2202/3481, -1/59x + 399/3481, 1/59x^2 + 73/3481x + 289/3481],$$

 $b = [0, 3481/2202x^2 - 13924/1101x + 10443/734, -3481/2202x^3 - 6962/1101x + 3481/2202].$ 

Since b[1] = 0 the loop terminates. We have LC(a[1]) = 2202/3481 and thus normalize to

$$a = [x^2 - 1, -59/2202x + 133/734, 59/2202x^2 + 73/2202x + 289/2202].$$

We check that indeed  $x^2 - 1 = (-59/2202x + 133/734)(x^5 + 3x^3 - x^2 - 4x + 1) + (59/2202x^2 + 73/2202x + 289/2202)(x^4 - 8x^3 + 8x^2 + 8x - 9).$ 

Here is a formal statement of the Rabin test:

## Lemma 3 (Rabin test)

The polynomial  $f(x) \in \mathbb{F}_q[x]$  of degree  $\deg(f) = m$  is irreducible if and only if

$$f(x)|x^{q^m} - x$$

and for all primes d < m dividing m one has

$$\gcd(f(x), x^{q^a} - x) = 1.$$