

**TECHNISCHE UNIVERSITEIT EINDHOVEN**  
**Faculty of Mathematics and Computer Science**  
**Exam Cryptology, Tuesday 31 October 2017**

Name :

TU/e student number :

Exercise	1	2	3	4	5	6	total
points							

**Notes:** Please hand in *this sheet* at the end of the exam. You may keep the sheet with the exercises.

This exam consists of 6 exercises. You have from 13:30 – 16:30 to solve them. You can reach 100 points.

Make sure to justify your answers in detail and to give clear arguments. Document all steps, in particular of algorithms; it is not sufficient to state the correct result without the explanation. If the problem requires usage of a particular algorithm other solutions will not be accepted even if they give the correct result.

All answers must be submitted on TU/e letterhead; should you require more sheets ask the proctor. State your name on every sheet.

Do not write in red or with a pencil.

You are allowed to use any books and notes, e.g. your homework. You are not allowed to use the textbooks of your colleagues.

You are allowed to use a calculator without networking abilities. Usage of laptops and cell phones is forbidden.



1. This problem is about the Diffie-Hellman key exchange. The system uses the multiplicative group  $\mathbb{F}_p^*$  modulo the prime  $p = 23689$ . The element  $g = 11 \in \mathbb{F}_{23689}^*$  has order 23688 and is thus a generator of the full multiplicative group.
  - (a) Alice chooses  $a = 222$  as her secret key. Compute Alice's public key. 1 point
  - (b) Alice receives  $h_b = g^b = 22938$  from Bob as his Diffie-Hellman keyshare.  
Compute the key shared between Alice and Bob, using Alice's secret key  $g^a$  from the first part of this exercise. 2 points
  
2. This problem is about RSA encryption.
  - (a) Alice chooses  $p = 439$  and  $q = 349$ . Compute Alice's public key  $(n, e)$ , using  $e = 2^{16} + 1$ , and the matching private key  $d$ . 2 points
  - (b) Bob uses public key  $(n, e) = (443507, 11)$  and secret key  $d = 241187$ . He receives ciphertext  $c = 64649$ .  
Decrypt the ciphertext. 2 points
  - (c) Decrypt the same message as under b) but this time using RSA with CRT for  $p = 659$  and  $q = 673$ . Make sure to document your computation, i.e., state the values for  $c_p, d_p, \dots$  4 points
  
3. This exercise is about computing discrete logarithms in the multiplicative group of  $\mathbb{F}_p$  for  $p = 23689$ . The element  $g = 11$  has order  $\ell = 23688$ . The factorization of  $p - 1$  is  $p - 1 = 2^3 \cdot 3^2 \cdot 7 \cdot 47$ . Use the Pohlig-Hellman attack to compute the discrete logarithm  $b$  of Bob's key  $h_b = g^b = 22938$ , i.e.
  - (a) Compute  $b$  modulo  $2^3$  by first computing  $b$  modulo 2, then modulo  $2^2$  and finally modulo  $2^3$ . 4 points
  - (b) Compute  $b$  modulo  $3^2$  by first computing  $b$  modulo 3 and then modulo  $3^2$ . 5 points
  - (c) Compute  $b$  modulo 7. 4 points
  - (d) Compute  $b$  modulo 47 using the Baby-Step Giant-Step attack in the subgroup of order 47 7 points
  - (e) Combine the results above to compute  $b$ .  
Verify your answer. 4 points

4. This exercise is about factoring  $n = 443507$ .
- (a) Use the  $p - 1$  method to factor  $n = 443507$  with basis  $a = 13$  and exponent  $s = \text{lcm}\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13\}$ . Make sure to state the value for  $s$  and the result of the exponentiation modulo  $n$ . Determine both factors of  $n$ . 4 points
- (b) Use Pollard's rho method for factorization to find a factor of 329 with iteration function  $x_{i+1} = x_i^2 + 3$  and Floyd's cycle finding method, i.e. after each increment in  $i$  compute  $\text{gcd}(x_{2i} - x_i, 329)$  until a non-trivial gcd is found. Start with  $x_0 = 3$ . 5 points
- (c) Use the result of b) to explain why the factorization in a) was successful. Note that  $673 - 1 = 2^5 \cdot 3 \cdot 7$  (factored completely) and  $659 - 1 = 2 \cdot 329$ . 3 points
5. (a) Find all affine points, i.e. points of the form  $(x, y)$ , on the Edwards curve
- $$x^2 + y^2 = 1 + 5x^2y^2$$
- over  $\mathbb{F}_{17}$ . 9 points
- (b) Verify that  $P = (5, 10)$  is on the curve. Compute the order of  $P$ .  
**Hint:** You may use information learned about the order of points on Edwards curves. 10 points
- (c) Translate the curve **and**  $P$  to Montgomery form
- $$Bv^2 = u^3 + Au^2 + u,$$
- i.e. compute  $A, B$  and the resulting point  $P'$ .  
 Verify that the resulting point  $P'$  is on the Montgomery curve. 6 points
- (d) The point  $Q = (1, 16)$  is on the Montgomery curve with  $A = 14, B = -1$  over  $\mathbb{F}_{17}$ . Compute  $3Q$ . 10 points
6. Lots of applications in cryptography require random numbers. The *power generator* generates random numbers in  $\mathbb{F}_p^*$  by taking random powers of a generator, i.e., computing random number  $x$  as  $x = g^r$  in  $\mathbb{F}_p$  for some fixed  $g$ .

- (a) Company C wants to generate numbers coprime to 3, 5, 7, 11, and 13. They choose to pick 5 small random numbers  $r_1, r_2, \dots, r_5$ , compute

$$x \equiv 2^{r_1} \pmod{3}$$

$$x \equiv 3^{r_2} \pmod{5}$$

$$x \equiv 3^{r_3} \pmod{7}$$

$$x \equiv 2^{r_4} \pmod{11}$$

$$x \equiv 2^{r_5} \pmod{13}$$

and then combine these five congruences using the Chinese Remainder Theorem (CRT) to a single number  $x$  modulo  $3 \cdot 5 \cdot 7 \cdot 11 \cdot 13 = 15015$ .

Explain why the resulting numbers are coprime to 3, 5, 7, 11, and 13.

Compute how many different numbers can be generated using this method.

6 points
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- (b) Company S wants to simplify the code and picks a single number as generator, so  $x$  is computed picking 5 small random numbers  $r_1, r_2, \dots, r_5$  as before and solving the following CRT for  $x$ .

$$x \equiv 2^{r_1} \pmod{3}$$

$$x \equiv 2^{r_2} \pmod{5}$$

$$x \equiv 2^{r_3} \pmod{7}$$

$$x \equiv 2^{r_4} \pmod{11}$$

$$x \equiv 2^{r_5} \pmod{13}$$

Compute how many different numbers can be generated using the method of company S.

3 points
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- (c) Impatient company I additionally wants to avoid the CRT step and generates numbers coprime to 15015 by taking a larger random number  $r < 15015$  and computing

$$x \equiv 5477^r \pmod{15015}.$$

Compute how many different numbers can be generated using the method of company I.

Verify your answer.

9 points
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