Authenticated Encryption Requirements

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Directions in Authenticated Ciphers, 2012

Many desirable attributes

- High security
- Computationally cheap
- Low latency
- Compact in software and/or hardware
- Re-use existing cryptographic components
- Randomized (no nonce)
- Misuse resistance
- Side channel resistance
- Forward security
- Postquantum
- Key agility
- Beyond birthday bound security
- Message length hiding

Domains of use

	message size	data rates	goals
Links	40 to	0.6 to	low latency
	2000 bytes	100 Gbit	
Internet	40 to	1 to 10 Mbit	
	2000 bytes		
Low power	1 to 100 bytes	20 to	low expansion
wireless		250 Kbits	compact
Data	512 to	400 Mbit	randomized?
at rest	4096 bytes		

AEAD in standards

AES-CCM 802.11i, 802.15, ESP, TLS protocols

AES-GCM 802.1AE (MACsec), INCITS Fibre Channel
(FC-SP), IKE, ESP, TLS, SSH, and SRTP,
P1619.1 and LTO-4 tape storage; Suite B

AES-OCB 802.11i

Camellia-GCM TLS

ARIA-GCM TLS

SEED-GCM TLS

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- Compact software implementations difficult
- Nonce re-use, short tags

$$Y_0 = \begin{cases} IV || 0^{31} 1 & \text{if len}(IV) = w - 32 \\ \text{GHASH}(H, \{\}, IV) & \text{otherwise.} \end{cases}$$

$$Y_i = \text{incr}(Y_{i-1}) \text{ for } i = 1, \dots, n$$

$$C_i = P_i \oplus E(K, Y_i) \text{ for } i = 1, \dots, n - 1$$

$$C_n^* = P_n^* \oplus \text{MSB}_u(E(K, Y_n))$$

$$T = \text{MSB}_t(\text{GHASH}(H, A, C) \oplus E(K, Y_0))$$

$$H = E(K, 0^w)$$

GCM evolution?

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Broadens applicability, but may not address all domains

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 - Available royalty-free worldwide